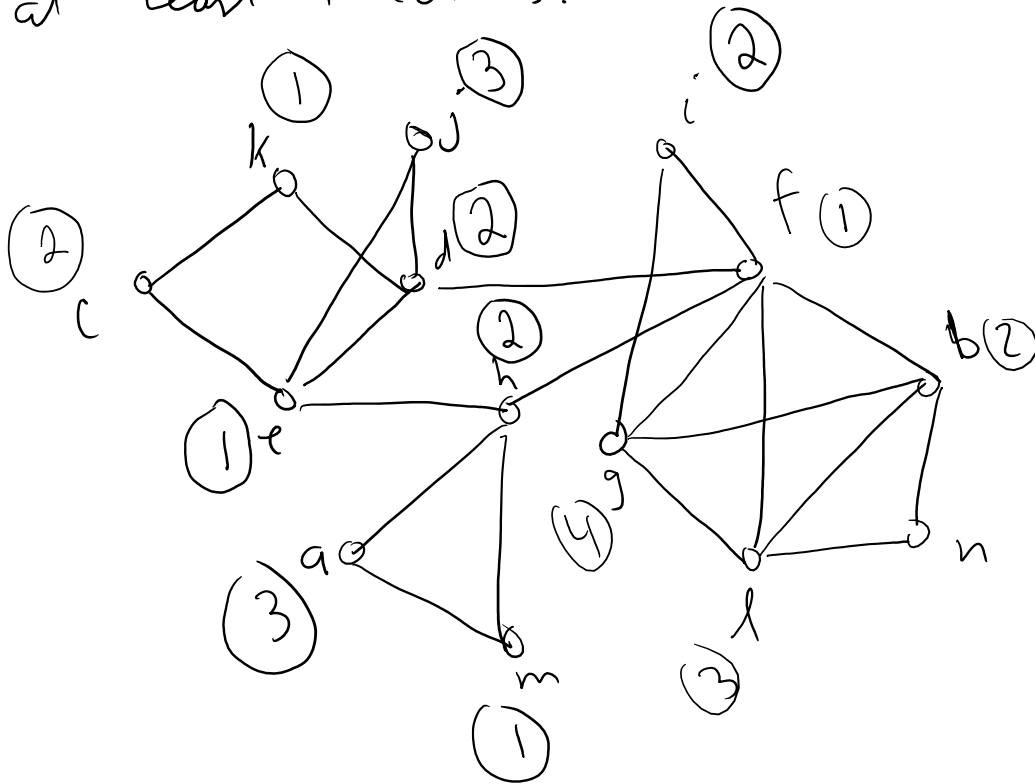


Homework 5 Solutions

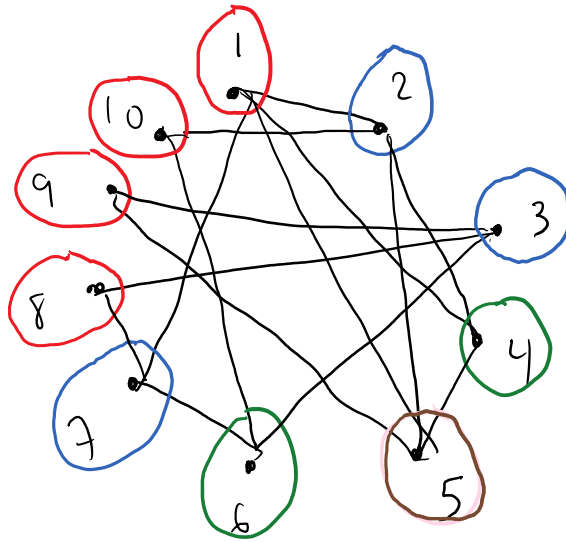
(14)

Since f, b, l, g forms a K_4 , we need at least 4 colors.



4 colors works, so $\chi(G) = 4$
(I used ①, ②, ③, ④ to represent colors)

(15)



We want to find the chromatic number (since you could think of a color as a storeroom).

Because of 1, 2, 4, 5 we need at least 4 colors.

Using red, blue, green and brown we show 4 colors suffice. Therefore we need 4 storerooms (and no more, no less).

(17) It's 2. You need at least 2 because the tree is connected, so it has at least one edge (which requires 2 colors). Since trees contain

no cycles, they contain no odd cycles, so they are 2-colorable. So $\chi(T) = 2$.

(21)

$$n_3 = 5.$$

$$n_{t+1} = \binom{t(n_t - 1) + 1}{n_t} n_t + t(n_t - 1) + 1$$

For each n_t -subset of \mathbb{I} , you make a copy of G_t
 $|\mathbb{I}| = t(n_t - 1) + 1$. So we have $\binom{|\mathbb{I}|}{n_t} n_t + |\mathbb{I}|$.

For example $n_4 = 5 \binom{13}{5} + 13 = 6448$

$$n_5 = 6448 \binom{25789}{6448} + 25789 \approx 1.2331 \times 10^{6300}$$

(22)

$$n_3 = 5$$

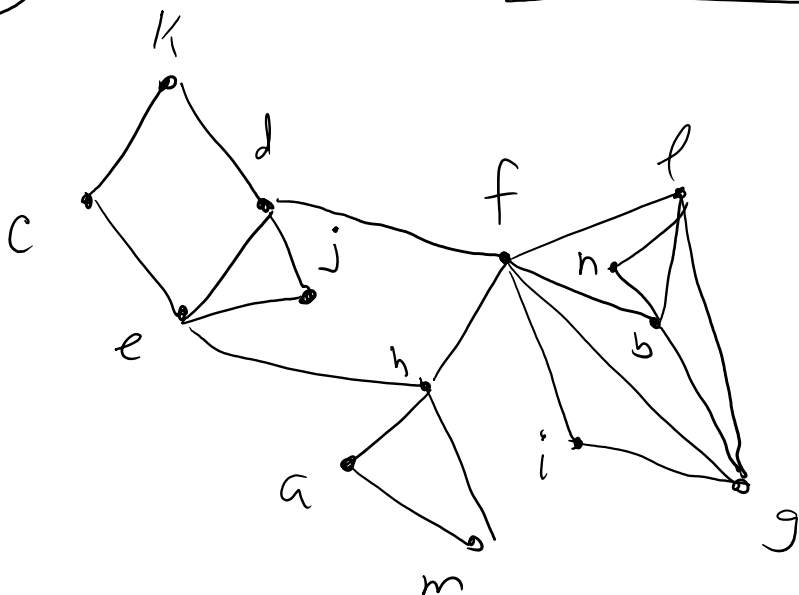
$$n_{t+1} = 2n_t + 1$$

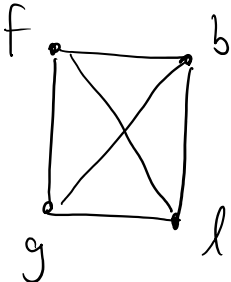
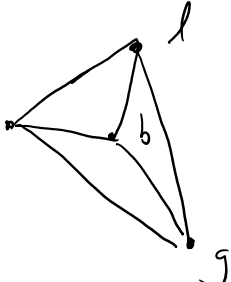
so $n_4 = 11, n_5 = 23, n_6 = 47, n_7 = 95, \dots$

$\tau + 1 \leq n$

28

G is planar



(one useful trick is to transform
 into )

32

Suppose every vertex is incident to at least 6 edges. Therefore every vertex has degree ≥ 6 .

$$\text{Then } |E| = \frac{1}{2} \sum_{v \in V} \deg(v) \geq 3n$$

But by Theorem 5.33, planar graphs have

at most $3n-6$ edges. Contradiction!

Therefore one vertex has degree at most 5.

(33) Cannot be the degree sequence of any graph:

c) $n=7$ $(5,4,4,3,2,1,0)$

The sum of the degrees is odd, which is impossible in a graph.

Two sequences that could be planar graphs:

Note $|E| \leq 3n-6$ so $\sum \deg(v) \leq 6n-12$.

e) has sum of degrees = $18 \leq 6(6)-12=24$ ✓

a) has sum = $18 \leq 6(6)-12=48$ ✓

So $\boxed{a, e}$

The one sequence that could be a tree:

Trees are planar so it's either a or e.

e can't be a tree. one vertex is connected to

all others, so if any other vertex has degree ≥ 2

then we get a cycle. Therefore the only

one that can be a tree is (a).

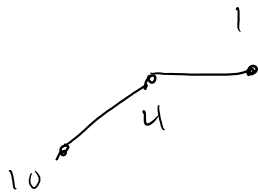
The one sequence that is Eulerian

(b) It's the only one where all vertices have even degree.

The one sequence that must be Hamiltonian

(d) Every vertex has degree $\geq \lceil \frac{10}{2} \rceil = 5$.

(37)



First leaf is 2 so 1

Next leaf is 3 connected to 4 1 4

Next leaf is 5 " " 6 1 4 6

Next 6 " " 9 1 4 6 9

7 → 4 1 4 6 9 4

8 → 9 1 4 6 9 4 9

9 → 1 1 4 6 9 4 9 1

1 → 4 1 4 6 9 4 9 1 4

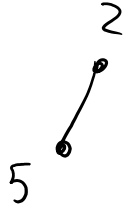
So

1 4 6 9 4 9 1 4

41 The tree is from $\{1, 2, 3, 4, 5, 6, 7\}$

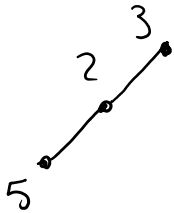
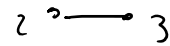
The initial leaves are 5, 6, 7 code
2 3 1 3 4

Start.



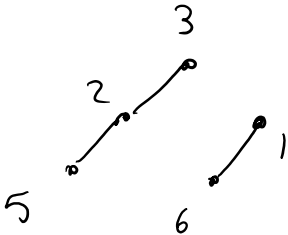
now leaves $\{6, 7, 2\}$

2 is smallest so



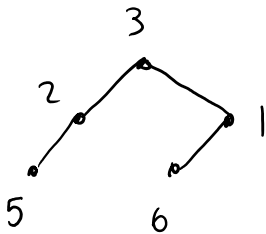
leaves $\{6, 7\}$

6 is smallest so $6 \rightarrow 1$



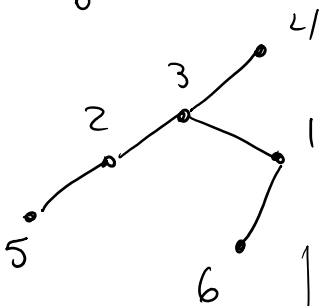
leaves $\{7, 1\}$

1 smallest so $1 \rightarrow 3$



leaves $\{7, 3\}$

3 smallest so $3 \rightarrow 4$



leaves $\{7, 4\}$

so $4 \rightarrow 7$

